# Introduction to IPv6 

## Scalable Infrastructure Workshop AfNOG

## Agenda

- Background
- Protocols \& Standards
- Addressing

ㅁ Co-existence \& Transition

## Early Internet History

- Late 1980s
- Exponential growth of the Internet
- Late 1990: CLNS proposed as IP replacement
- 1991-1992
- Running out of "class-B" network numbers
- Explosive growth of the "default-free" routing table
- Eventual exhaustion of 32-bit address space

ㅁ Two efforts - short-term vs. long-term

- More at "The Long and Windy ROAD"
http://rms46.vlsm.org/1/42.html


## Early Internet History

- CIDR and Supernetting proposed in 1992-3
- Deployment started in 1994
- IETF "ipng" solicitation - RFC1550, Dec 1993
- Direction and technical criteria for ipng choice RFC1719 and RFC1726, Dec 1994
- Proliferation of proposals:
- TUBA - RFC1347, June 1992
- PIP - RFC1621, RFC1622, May 1994
- CATNIP - RFC1707, October 1994
- SIP - RFC1710, October 1994
- NIMROD - RFC1753, December 1994
- ENCAPS - RFC1955, June 1996


## Early Internet History

$\square$ Other activities included:

- Development of NAT, PPP, DHCP,...
- Some IPv4 address reclamation
- The RIR system was introduced
$\square \rightarrow$ Brakes were put on IPv4 address consumption
- IPv4 32 bit address $=4$ billion hosts
- HD Ratio (RFC3194) realistically limits IPv4 to 250 million hosts


## Recent Internet History

The "boom" years $\rightarrow 2001$

- IPv6 Development in full swing
- Rapid IPv4 consumption
- IPv6 specifications sorted out
- (Many) Transition mechanisms developed
- 6bone
- Experimental IPv6 backbone sitting on top of Internet
- Participants from over 100 countries
- Early adopters
- Japan, Germany, France, UK,...


## Recent Internet History

The "bust" years: $2001 \rightarrow 2004$

- The DotCom "crash"
- i.e. Internet became mainstream
$\square$ IPv4:
- Consumption slowed
- Address space pressure "reduced"
- Indifference
- Early adopters surging onwards
- Sceptics more sceptical
- Yet more transition mechanisms developed


## $2004 \rightarrow 2011$

- Resurgence in demand for IPv4 address space
- IANA pool running low, several /8 block given to RIRs each year
- February 2011 final allocation of /8 blocks from IANA to RIRs


## 2012

 $\rightarrow$ NowSeveral RIRs in special policies for final block

- APNIC, RIPE NCC, ARIN
- ISPs can't get IPs as they need
- ISPs don't give public IPs to customer
- $\rightarrow$ that customer can't be directly reached from other networks
- Market for IPv4 addresses:
- Creates barrier to entry
- Condemns the less affluent to use of NATs
- IPv6 offers vast address space
- The only compelling reason for IPv6


## Current Situation

- IPv6 is being deployed
- It is the only sustainable plan forward
- Private sector requires a business case to "migrate"
$\square$ No easy Return on Investment (RoI) computation
- Some measure over 6\% of traffic is IPv6
- Network operators are at very different stages of deploying IPv6
- the have their head still in the sand
- Something needs to be done to sustain the Internet growth
- IPv6 or NAT or both or something else?


## Do we really need a larger address

- Internet population
- ~630 million users end of 2002 - 10\% of world pop.
- ~1320 million users end of 2007 - 20\% of world pop.
- Future? (World pop. ~9B in 2050)

ㅁ US uses $90 / 8 \mathrm{~s}$ - this is 6.4 IPv4 addresses per person

- Repeat this the world over...
- 6 billion population could require 26 billion IPv4 addresses
- (7 times larger than the IPv4 address pool)

ㅁ Emerging Internet economies need address space:

- China uses more than 249 million IPv4 addresses today (14.8/8s)
- Source: http://resources.potaroo.net/iso3166/v4cc.html


## Do we really need a larger address

- RFC 1918 is not sufficient for large environments
- Cable Operators (e.g. Comcast - NANOG37 presentation)
- Mobile providers (fixed/mobile convergence)
- Large enterprises
- The Policy Development process of the RIRs turned down a request to increase private address space
- RIR membership guideline is to use global addresses instead
- This leads to an accelerated depletion of the global address space
- Some want 240/4 as new private address space
- But how to back fit onto all TCP/IP stacks released since 1995?


## Do we really need a larger address

space?

- Large variety of proposals to "make IPv4 last longer" to help with IPv6 deployment
- NAT444
- Lots of IPv4 NAT
- Dual Stack Lite
- Improvement on NAT464
- Activity of IETF Softwires Working Group
- NAT64 \& IVI
$\square$ Translation between IPv6 and IPv4
$\square$ Activity of IETF Behave Working Group
- 6rd
$\square$ Dynamic IPv6 tunnel from SP to customer


## IPv6 OS and Application Support

- All software vendors officially support IPv6 in their latest Operating System releases
- Application Support
- Applications must be IPv4 and IPv6 agnostic
- User should not have to "pick a protocol"
- Successful deployment is driven by Applications


## ISP Deployment Activities

- Several Market segments
- IX, Carriers, Regional ISP, Wireless
- ISP have to get an IPv6 prefix from their Regional Registry
- Most carriers have deployed IPv6 in their core:
- And do offer IPv6 services to Internet customers
- Dual stack
- Several "transition mechanisms" to work around equipment restrictions
- Mostly for the access network


## Why not use Network Address

Translation?

- Private address space and Network address translation (NAT) could be used instead of IPv6
- But NAT has many serious issues:
- Breaks the end-to-end model of IP
- Breaks end-to-end network security
- Serious consequences for Lawful Intercept
- Non-NAT friendly applications means NAT has to be upgraded
- Some applications don't work through NATs
- Layered NAT devices
- Mandates that the network keeps the state of the connections
- How to scale NAT performance for large networks??
- Makes fast rerouting and multihoming difficult
- How to offer content from behind a NAT?


## Conclusion

- There is a need for a larger address space
- IPv6 offers this - will eventually replace NAT
- But NAT will be around for a while too
- Market for IPv4 addresses exists
- IPv6 deployment will reduce dependency on IPv4
$\square$ Many challenges ahead
- It's part of our work


# Protocols \& Standards 

## So what has really changed?

- Expanded address space
- Address length quadrupled to 16 bytes
- Header Format Simplification
- Fixed length, optional headers are daisy-chained
- IPv6 header is twice as long (40 bytes) as IPv4 header without options (20 bytes)
- No checksum at the IP network layer
- No hop-by-hop segmentation
- Path MTU discovery
- 64 bits aligned
- Authentication and Privacy Capabilities
- IPsec is mandated
- No more broadcast


## IPv4 and IPv6 Header Comparison

## IPv4 Header

| Version | IHL | Type of <br> Service |  | Total Length |
| ---: | :--- | :--- | :--- | :--- |
| Identification | Flags | Fragment <br> Offset |  |  |
| Time to Live | Protocol | Header Checksum |  |  |
|  |  |  |  |  |

Destination Address

## Options

Field's name kept from IPv4 to IPv6

Name and position changed in IPv6

Fields not kept in IPv6

New field in IPv6
Legend

## Padding

IPv6 Header

| Version | Flow Laffic Class |  |
| :---: | :---: | :---: |
| Payload Length | Next Header | Hop Limit |
| Source Address |  |  |

## Larger Address Space



- IPv4
- 32 bits
- = 4,294,967,296 possible addressable devices
- IPv6
$\square 128$ bits: 4 times the size in bits
ㅁ $=3.4 \times 10^{38}$ possible addressable devices
$\square=340,282,366,920,938,463,463,374,607,431,768,211,456$


## How was the IPv6 Address Size

 Chosen?- Some wanted fixed-length, 64-bit addresses
- Easily good for 1012 sites, 1015 nodes, at . 0001 allocation efficiency (3 orders of magnitude more than IPv6 requirement)
- Minimizes growth of per-packet header overhead
- Efficient for software processing
- Some wanted variable-length, up to 160 bits
- Compatible with OSI NSAP addressing plans
- Big enough for auto-configuration using IEEE 802 addresses
- Could start with addresses shorter than 64 bits \& grow later
- Settled on fixed-length, 128-bit addresses


## IPv6 Address Representation

- 16 bit fields in case insensitive colon hexadecimal representation
- 2031:0000:130F:0000:0000:09C0:876A:130B
- Leading zeros in a field are optional:
- 2031:0:130F:0:0:9C0:876A:130B
- Successive fields of 0 represented as ::, but only once in an address:
- 2031:0:130F::9C0:876A:130B is ok
- 2031::130F::9C0:876A:130B is NOT ok

- 0:0:0:0:0:0:0:1 $\rightarrow$ ::1 (loopback address)
- 0:0:0:0:0:0:0:0 $\rightarrow$ :: (unspecified address)


## IPv6 Address Representation

- In a URL, it is enclosed in brackets (RFC3986)
- http://[2001:db8:4f3a::206:ae14]:8080/index.html
- Cumbersome for users
- Mostly for diagnostic purposes
- Use fully qualified domain names (FQDN)
- Prefix Representation
- Representation of prefix is same as for IPv4 CIDR
- Address and then prefix length
- IPv4 address:
- 198.10.0.0/16
- IPv6 address:
- 2001:db8:1200::/40


## IPv6 Addressing

- IPv6 Addressing rules are covered by multiples RFCs
- Architecture defined by RFC 4291
- Address Types are :
- Unicast : One to One (Global, Unique Local, Link local)
- Anycast : One to Nearest (Allocated from Unicast)
- Multicast : One to Many
- A single interface may be assigned multiple IPv6 addresses of any type (unicast, anycast, multicast)
- No Broadcast Address $\rightarrow$ Use Multicast


## IPv6 Addressing

| Type | Binary | Hex |
| :---: | :--- | :---: |
| Unspecified | $000 \ldots 0$ | $:: / 128$ |
| Loopback | $000 \ldots 1$ | $:: 1 / 128$ |
| Global Unicast Address | $001 \ldots$ | $2000:: / 3$ |
| Link Local Unicast Address | 1111111010 | FE80::/10 |
| Unique Local Unicast Address | 11111100 | 11111101 |
| Multicast Address | 11111111 | FF00::/7 |
| FF00::/8 |  |  |

## IPv6 Global Unicast Addresses



## Global Routing Prefix $\quad$ Subnet-id

## Interface ID

- IPv6 Global Unicast addresses are:
- Addresses for generic use of IPv6
- Hierarchical structure intended to simplify aggregation


## IPv6 Address Allocation



- The allocation process is:
- The IANA is allocating out of 2000::/3 for initial IPv6 unicast use
- Each registry gets a /12 prefix from the IANA
- Registry allocates a /32 prefix (or larger) to an ISP
- Policy is that an ISP allocates a /48 prefix to each end customer


## IPv6 Addressing Scope

- 64 bits reserved for the interface ID
- Possibility of 264 hosts on one network LAN
- Arrangement to accommodate MAC addresses within the IPv6 address
- 16 bits reserved for the end site
- Possibility of $2{ }^{16}$ networks at each end-site
- 65536 subnets equivalent to a /12 in IPv4 (assuming 16 hosts per IPv4 subnet)


## IPv6 Addressing Scope

$\square 16$ bits reserved for the service provider

- Possibility of $2{ }^{26}$ end-sites per service provider
- 65536 possible customers: equivalent to each service provider receiving a $/ 8$ in IPv4 (assuming a /24 address block per customer)
- 29 bits reserved for service providers
- Possibility of 229 service providers
- i.e. 500 million discrete service provider networks
$\square$ Although some service providers already are justifying more than a /32
- Equivalent to an eighth of the entire IPv4 address space


## Aggregation hopes



- Larger address space enables aggregation of prefixes announced in the global routing table
- Idea was to allow efficient and scalable routing
- But current Internet multihoming solution breaks this model


## Interface IDs

- Lowest order 64-bit field of unicast address may be assigned in several different ways:
- Auto-configured from a 64-bit EUI-64, or expanded from a 48-bit MAC address (e.g., Ethernet address)
- Auto-generated pseudo-random number (to address privacy concerns)
- Assigned via DHCP
- Manually configured


## EUI-64

Ethernet MAC address (48 bits)


64 bits version

Uniqueness of the MAC

Eui-64 address

| 00 90 27 FF FE 17 FC $0 F$ |
| :--- |
| $000000 \times 0$ |
| 02 90 27 |

- EUI-64 address is formed by inserting FFFE and OR'ing a bit identifying the uniqueness of the MAC address


## IPv6 Address Privacy (RFC 3041)

| 132 |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
| 2001 | 0db8 |  |  | Interface ID |

- Temporary addresses for IPv6 host client application, e.g. Web browser
- Intended to inhibit device/user tracking but is also a potential issue
- More difficult to scan all IP addresses on a subnet
- But port scan is identical when an address is known
- Random 64 bit interface ID, run DAD before using it
- Rate of change based on local policy
- Implemented initially on Microsoft Windows XP/Vista/7 only
- Other OSes have now also adopted this
- Can be activated on FreeBSD/Linux/MacOS with a system call


## IPv6 Auto-Configuration

- Stateless (RFC4862)
- Host autonomously configures its own Link-Local address
- Router solicitation are sent by booting nodes to request RAs for configuring the interfaces.
- Stateful
- DHCPv6 - required by most enterprises
- Renumbering
- Hosts renumbering is done by modifying the RA to announce the old prefix with a short lifetime and the new prefix
- Router renumbering protocol (RFC 2894), to allow domain-interior routers to learn of prefix introduction / withdrawal



## Auto-configuration

Mac address: 00:2c:04:00:FE:56

Host autoconfigured address is: prefix received + linklayer address

information (prefix, default route, ...)

- PC sends router solicitation (RS) message
- Router responds with router advertisement (RA)
- This includes prefix and default route
- PC configures its IPv6 address by concatenating prefix received with its EUI-64 address


## Renumbering

Mac address: 00:2c:04:00:FE:56


> Sends $N E W$ network-type information (prefix, default route, ...)

Host auto-configured address is:
NEW prefix received + SAME link-layer address


- Router sends router advertisement (RA)
- This includes the new prefix and default route (and remaining lifetime of the old address)
- PC configures a new IPv6 address by concatenating prefix received with its EUI-64 address
- Attaches lifetime to old address


## Unique-Local

## 128 Bits



## 7 Bits

- Unique-Local Addresses Used For:
- Local communications
- Inter-site VPNs
- Site Network Management systems connectivity
- Not routable on the Internet

ㅁ Reinvention of the deprecated site-local?

## Link-Local

## Remaining 54 Bits <br> Interface ID

1111111010

## FE80::/10

## 10 Bits

- Link-Local Addresses Used For:
- Communication between two IPv6 device (like ARP but at Layer 3)
- Next-Hop calculation in Routing Protocols
- Automatically assigned by Router as soon as IPv6 is enabled
- Mandatory Address
- Only Link Specific scope
- Remaining 54 bits could be Zero or any manual configured value


## Multicast use

- Broadcasts in IPv4
- Interrupts all devices on the LAN even if the intent of the request was for a subset
- Can completely swamp the network ("broadcast storm")
- Broadcasts in IPv6
- Are not used and replaced by multicast
- Multicast
- Enables the efficient use of the network
- Multicast address range is much larger


## IPv6 Multicast Address

- IP multicast address has a prefix FF00::/8
$\square$ The second octet defines the lifetime and scope of the multicast address.

| 8-bit | 4-bit | 4-bit | 112 -bit |
| :---: | :---: | :---: | :---: |
| 11111111 | Lifetime | Scope | Group-ID |


| Lifetime |  |
| :---: | :---: |
| 0 | If Permanent |
| 1 | If Temporary |


| Scope |  |
| :---: | :---: |
| 1 | Node |
| 2 | Link |
| 5 | Site |
| 8 | Organization |
| E | Global |

## IPv6 Multicast Address Examples

- RIPng
- The multicast address AllRIPRouters is FF02::9
$\square$ Note that 02 means that this is a permanent address and has link scope
- OSPFv3
- The multicast address AllSPFRouters is FF02::5
- The multicast address AllDRouters is FF02::6
- EIGRP
- The multicast address AllEIGRPRouters is FF02::A


## IPv6 Anycast

- An IPv6 anycast address is an identifier for a set of interfaces (typically belonging to different nodes)
- A packet sent to an anycast address is delivered to one of the interfaces identified by that address (the "nearest" one, according to the routing protocol's measure of distance).
- RFC4291 describes IPv6 Anycast in more detail
- In reality there is no known implementation of IPv6 Anycast as per the RFC
- Most operators have chosen to use IPv4 style anycast instead


## Anycast on the Internet

- A global unicast address is assigned to all nodes which need to respond to a service being offered
- This address is routed as part of its parent address block
$\square$ The responding node is the one which is closest to the requesting node according to the routing protocol
- Each anycast node looks identical to the other
- Applicable within an ASN, or globally across the Internet
- Typical (IPv4) examples today include:
- Root DNS and ccTLD/gTLD nameservers
- SMTP relays within ISP autonomous systems


## MTU Issues

- Minimum link MTU for IPv6 is 1280 octets (versus 68 octets for IPv4)
- $\Rightarrow$ on links with MTU < 1280, link-specific fragmentation and reassembly must be used
ㅁ Implementations are expected to perform path MTU discovery to send packets bigger than 1280
- Minimal implementation can omit PMTU discovery as long as all packets kept $\leq 1280$ octets
- A Hop-by-Hop Option supports transmission of "jumbograms" with up to 232 octets of payload


## Neighbour Discovery (RFCs 2461 \&

- Protocol built on top of ICMPv6 (RFC 4443)
- combination of IPv4 protocols (ARP, ICMP, IGMP,...)
- Fully dynamic, interactive between Hosts \& Routers
- defines 5 ICMPv6 packet types:
$\square$ Router Solicitation / Router Advertisements
$\square$ Neighbour Solicitation / Neighbour Advertisements
$\square$ Redirect


## IPv6 and DNS



## IPv6 Technology Scope

 IP Service IPver solution IPvo solutionAddressing Range

Autoconfiguration

Security

DHCP

IPSec

128-bit, Multiple Scopes

Serverless, Reconfiguration, DHCP

IPSec Mandated, works End-to-End

Mobile IP with Direct Routing

Mobile IP
Differentiated Service, Integrated Service

Differentiated Service, Integrated Service

IP Multicast
IGMP/PIM/Multicast BGP

MLD/PIM/Multicast BGP,Scope Identifier

## What does IPv6 do for:

- Security
- Nothing IPv4 doesn't do - IPSec runs in both
- But IPv6 architecture mandates IPSec
- QoS
- Nothing IPv4 doesn't do -
$\square$ Differentiated and Integrated Services run in both
- So far, Flow label has no real use


## IPv6 Status - Standardisation

ㅁ Several key components on standards track...

Specification (RFC2460)
ICMPv6 (RFC4443)
RIP (RFC2080)
IGMPv6 (RFC2710)
Router Alert (RFC2711)
Autoconfiguration (RFC4862)
DHCPv6 (RFC3315 \& 4361)
IPv6 Mobility (RFC3775)
GRE Tunnelling (RFC2473)
DAD for IPv6 (RFC4429)
ISIS for IPv6 (RFC5308)

- IPv6 available over:

PPP (RFC5072)
FDDI (RFC2467)
NBMA (RFC2491)
Frame Relay (RFC2590)
IEEE1394 (RFC3146)
Facebook (RFC5514)

Neighbour Discovery (RFC4861 \& 4311)
IPv6 Addresses (RFC4291 \& 3587)
BGP (RFC2545)
OSPF (RFC5340)
Jumbograms (RFC2675)
Radius (RFC3162)
Flow Label (RFC3697)
Mobile IPv6 MIB (RFC4295)
Unique Local IPv6 Addresses (RFC4193)
Teredo (RFC4380)

Ethernet (RFC2464)
Token Ring (RFC2470)
ATM (RFC2492)
ARCnet (RFC2497)
FibreChannel (RFC4338)

# Addressing 

## Getting IPv6 address space

- Become a member of your Regional Internet Registry and get your own allocation
- Require a plan for a year ahead
- General allocation policies and specific details for IPv6 are on the individual RIR website or
ㅁ Take part of upstream ISP's PA space or
- Use 6to4 (absolutely last resort)
- There is plenty of IPv6 address space
- The RIRs require high quality documentation


## Getting IPv6 address space

- From the RIR
- Receive a /32 (or larger if you have more than 65k / 48 assignments)
- From your upstream ISP
- Get one /48 from your upstream ISP
- More than one /48 if you have more than 65k subnets
- Use 6to4 (not recommended)
- Take a single public IPv4 /32 address
- 2002: <ipv4 /32 address>::/48 becomes your IPv6 address block, giving 65k subnets
- Requires a 6to4 gateway
- Routing/performance can be "strange"


## Addressing Plans - ISP

Infrastructure

- ISPs should receive /32 from their RIR
- Address block for router loop-back interfaces
- Generally number all loopbacks out of one /64
- Address block for infrastructure
- /48 allows 65k subnets
- /48 per PoP or region (for huge networks)
- /48 for whole backbone (commonly used by most ISPs)
- Summarise between sites if it makes sense


## Addressing Plans - ISP

 Infrastructure- What about LANs?
- /64 per LAN
- What about Point-to-Point links?
- Expectation is that /64 is used
- People have used /127s and /126s
- Mobile IPv6 Home Agent discovery won't work (doesn't matter on PtP links)
- People have used $/ 112 s$
- Leaves final 16 bits free for node IDs
- See RFC3627 for more discussion
- Discussion about /127 for PtP links: www.ietf.org/internet-drafts/draft-kohno-ipv6-prefixlen-p2p-01.txt


## Addressing Plans - Customer

- Customers get one /48
- Unless they have more than 65k subnets in which case they get a second /48 (and so on)
- (Still on going industry discussion about giving "small" customers a /56 or a /60 and single LAN end-sites a / 64)
- Should not be reserved or assigned on a per PoP basis
- ISP iBGP carries customer nets
- Aggregation within the iBGP not required and usually not desirable
- Aggregation in eBGP is very necessary


# Addressing Plans - ISP Infrastructure <br> ㅁ Phase One 



## Addressing Plans

Planning

- Registries will usually allocate the next block to be contiguous with the first allocation
- Minimum allocation is /32
- Very likely that subsequent allocation will make this up to a /31
- So plan accordingly


## Addressing Tools

ㅁ Examples of IP address tools (which support IPv6 too):

- IPAT
- ipv6gen
- sipcalc
- freeipdb
http://nethead.de/index.php/ipat
http://techie.devnull.cz/ipv6/ipv6gen/
http://www.routemeister.net/projects/sipcalc/
http://home.globalcrossing.net/~freeipdb/


# Transition \& Coexistence 

## IPv4-IPv6 Co-existence/Transition

- A wide range of techniques have been identified and implemented, basically falling into three categories:
- Dual-stack techniques, to allow IPv4 and IPv6 to coexist in the same devices and networks
- Tunneling techniques, to avoid dependencies when upgrading hosts, routers, or regions
- Translation techniques, to allow IPv6-only devices to communicate with IPv4-only devices
- Expect all of these to be used, in combination


## Dual Stack Approach



Data Link (Ethernet)


- Dual stack node means:
- Both IPv4 and IPv6 stacks enabled
- Applications can talk to both
- Choice of the IP version is based on name lookup and application preference


## Dual Stack \& DNS



2001:db8:1::1

- On a system running dual stack, an application that is both IPv4 and IPv6 enabled will:
- Ask the DNS for an IPv6 address (AAAA record)
- If that exists, IPv6 transport will be used
- If it does not exist, it will then ask the DNS for an IPv4 address (A record) and use IPv4 transport instead


## Using Tunnels for IPv6 Deployment

- Many techniques are available to establish a tunnel:
- Manually configured
- Manual Tunnel (RFC 4213) \& GRE (RFC 2473)
- Semi-automated
- Tunnel broker
- Automatic
- 6rd (RFC 5569) \& 6to4 (RFC 3056)


## ISATAP \& TEREDO are more useful

 for Enterprises than for ServiceProviders - but Security Concerns!!

- ISATAP (RFC 4214) \& TEREDO (RFC 4380)
- Opinion today is that any type of tunneling is "bad" and native is "good"


## Summary

- IPv6 offers vast address space
$\square$ Distinct addressing hierarchy between ISPs, end-sites, and LANs
- Planning is not so "confined" as for IPv4
- Coexistence with, NOT replacement of IPv4
- Clients prefer IPv6 before IPv4
- If IPv6 is configured \& available

